On equality of two classes of homomorphism-homogeneous relational structures

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Motivate

Definition (Ultrahomogeneity)

 \mathcal{A} is ultrahomogeneous if every isomorphism $f: \mathcal{A} \to \mathcal{B}$ between finite substructures $\mathcal{A}, \mathcal{B} \subset \mathcal{A}$ can be extended into automorphism



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How far structures from homogeneity

- ► Relational complexity r (Cherlin, Martin, Saracino 1996) i.e. expand language with relations of arity ≤ r s.t. autmomorphism group remains and this lift is homogeneous
- For graphs (H., Hubička, Nešetřil, 2015) used result of (Hubička, Nešetřil, 2014) to homogenize graphs with forbidden homomorphism.

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- For graphs (H., Hubička, Nešetřil, 2015) used result of (Hubička, Nešetřil, 2014) to homogenize graphs with forbidden homomorphism.
- Relax definition of homogeneity Use homomorphism instead of isomorphism It also have Fraïssé type results

(Ultra) Homogeneity of structures

Considered structures a

- ▶ relational structure $A = (A, \mathcal{R}_A)$ where $\mathcal{R}_A = (R_A^i; i \in I)$
- Usually interpreted as colored graphs (consider just binary relations - see later)

Classifications usually differs depending on

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Classification for finite graphs (Gardiner 1976)

Combinatorial argument utilizing finiteness of structures



Rado graph ${\mathcal R}$

Countably infinite random graph

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- Useful property

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(*) \forall X, Y finite \exists z \text{ s.t. } z \sim x \ \forall x \in X \text{ and } z \nsim y \ \forall y \in Y
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- Useful properties
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 - \blacktriangleright Universality: All finite graphs can be embedded into ${\cal R}$
 - Homogeneity: Graph with this property is homogeneous

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- Idea of proof(s)
 - Start with A, B finite and isomorphism $f : A \rightarrow B$
 - Iteratively construct one vertex extension of f using (*)
 - Automorphism is the union of partial maps

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Important notes

- It is, of course, a Fraïssé limit for class of all graphs
 - Showing this gives us all properties above
- Part of complete classification (Lachlan, Woodrow 1980)

Variants of homogeneity (Cameron, Nešetřil 2006)

- homomorphism-homogeneity (HH)
 - Iocal homomorphism \rightarrow homomorphism
- monomorphism-homogeneity (MH)
 - $\blacktriangleright \ \text{local monomorphism} \rightarrow \text{homomorphism}$

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Problems

- Classification beyond finite graphs
 - Finite HH Graphs (Cameron, Nešetřil 2006)
- $HH \subseteq MH$, HH = MH ?
 - ► For countable graphs YES! (Rusinov, Schweitzer 2010)
 - For general structures NO!

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 $\chi(v) \leq_P \chi(f(v))$ and $\xi(x,y) \leq_Q \xi(f(x),f(y))$

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Corresponding classes HH_{P,Q} and HH_Q

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- Determine borderline property for equality.
- Is it vertex coloring?



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A countable graph contains \mathcal{R} as a spanning subgraph if and only if it has the (\triangleright) property. Moreover any such graph is **HH** and **MH**.

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Structure \mathcal{R}_n

- ▶ Let C_n be a class of finite graphs with edges colored by F_n (F_n is antichain extended by minimal element 0)
- \mathcal{R}_n is universal for \mathcal{C}_n and homogeneous

Properties of \mathcal{R}_n

Let G be F_n -colored graph

 (\Diamond_n) Let $G_1, G_2, \ldots, G_{n+1}$ be finite disjoint subsets of G then there exists $x \in G \setminus G_{n+1}$ s.t. each vertex of $G_i \sim_i x$ for $1 \le i \le n$ and for each $y \in G_{n+1}$ a pair xy is non-edge.

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Using extension property (\Diamond_n) to find one vertex extension

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Lesson learned: Vertex coloring is not needed!

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- Q is directed set (upper bound for any pair)
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- Enumerate
 - all finite subsets of X as $\{Y_i : i \in \omega\}$
 - ▶ all functions t_j^i : $Y_i \to F_n$ there are $j \in \{2, \dots, (n+1)^{|Y_i|}\}$

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Note that

- Each point of construction uses only finitely many elements it's possible
- Condition 2. satisfied by construction and 1. follows

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 - \blacktriangleright \geq 1 non-edge between vertices of equal color remains
 - This contradicts HH

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 - ► This contradicts **HH**
 - extension property (\Diamond_n) imply **MH**

Lemma (Aranda, H., 2018)

Let P and Q be finite partially ordered sets. If $\mathbf{MH}_{P,Q} = \mathbf{HH}_{P,Q}$, then Q is a directed set.

Basic idea: Prove contrapositive: Q is not directed set

- Q has maximal elements R_1, \ldots, R_n
- ▶ |*P*| = *m*
- Let M₀ be Fraïssé limit of graphs colored by F_n
- ▶ Partition *M*⁰ as in previous lemma
 - Elements of *P* are e_1, \ldots, e_m and assign $e_i \rightarrow C_i$
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What about sufficient condition for equality of classes?

Construction of example structure M

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not extendable to $d \in M_c^0$

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Assume image of f to have exactly 2 elements

- ▶ from {a, b, c} say b, c
- study effect of f on {a, b, c} (case analysis)





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Let *P* and *Q* be finite partially ordered sets. $\mathbf{MH}_{P,Q} = \mathbf{HH}_{P,Q}$ iff *Q* is a linear order.

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 - Choose s^0 such that $\forall s \in S$ s.t. $\varphi_{u,s} \preceq \varphi_{u,s^{\bar{u}}}$
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 - use another realization of s^0 so $f \cup \{(u, (f \upharpoonright s^0)(u_0)\}\}$



Backward idea: If Q is not linear then $\exists M \text{ s.t. } M \text{ is } \mathbf{MH}_{P,Q}$ but not $\mathbf{HH}_{P,Q}$. We know that Q is finite directed set

• Let top element be 1 and let P_1, \ldots, P_n be top elements $Q \setminus 1$

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Construct M which is connected in $\mathbb{1}$ as well as P_i

Start with M₁ Rado graph in maximal color





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- Start with M₁ Rado graph in maximal color
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- Start with M₁ Rado graph in maximal color
- ▶ Partition M_1 into n! sets and fill all non-edges with colors from Q s.t. all are used (meeting condition from lemma)
- Create n + 1 such M_i



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- Start with M₁ Rado graph in maximal color
- ▶ Partition *M*¹ into *n*! sets and fill all non-edges with colors from *Q* s.t. all are used (meeting condition from lemma)
- Create n + 1 such M_i
- Add new vertices $x_1, x_2, \ldots, x_{n+1}$



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- Connect x_i to parititions of M_j s.t. x_i is connected in any combination of colors to all M_j, j ≠ i



Connect x_j and M_i

Coloring edges between x_j and M_i

• Indices of parts of M_i corresponds to $\sigma \in S_n$, i.e.

$$M = \bigcup_{\sigma \in S_n} M_i^{\sigma}$$





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For edges between x_i and M_i define color choosing function

$$c_i: \{x_1,\ldots,x_{i-1},x_{i+1},\ldots,x_n\} \rightarrow \{1,2,\ldots,n\}$$





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 M_i

• Color M_i^{σ} and x_j by $P_{\sigma(c_i(x_j))}$



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This construction generalize





 M_i

Extended example - properties

Backward idea - homomorphism-homogeneity of M: M is not $HH_{P,Q}$

- x_i and x_j are in distance 3 in $\mathbb{1}$
- Take distinct $j, k, \ell \in \{1, \dots, n+1\}$
- ► Local homomorphism $x_k \mapsto x_k, x_\ell \mapsto x_k, x_j \mapsto x_j$ cannot be extended
 - P_{i_1} and P_{i_2} has only $\mathbbm{1}$ above



Extended example - properties

Backward idea - homomorphism-homogeneity of M: M is not $HH_{P,Q}$



M is $\mathbf{MH}_{P,Q}$ Let $f: H \to K$ be surjective monomorphism

- If $|K \cap \{x_1, \ldots, x_{n+1}\}| \le 1$ we have infinitely many cones
- Assume K contains at least two vertices from $\{x_1, \ldots, x_{n+1}\}$
 - ▶ Note that *x_ix_j* is only non-edge in *M*, i.e.
 - ▶ its preimage is contained in {x₁,..., x_{n+1}}

Final claim finishing the proof

Claim

Given any $F \subset M \setminus \{x_1, \ldots, x_{n+1}\}$, $S \subseteq \{x_1, \ldots, x_{n+1}\}$ and injective $t : S \to \{\mathbb{1}, P_1, \ldots, P_n\}$, there $\exists v \in M$ that is connected to all of F by edges of type $\mathbb{1}$ and satisfies t over S, i.e., $\varphi_{v,S} = t$.



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For $v_i \in M_i$ and any F

▶ $v_i \sim_1 u$ for $u \in M_j \cup F$ for $j \neq i$ by construction of M

▶ $v_i \sim_1 u$ for $u \in M_i \cup F$ by Rado-ness of M_i

For injective t

• using 1 for x_i choose v from M_i connected correctly (correct σ)

others similarly

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We use this claim to proof that M is $\mathbf{MH}_{P,Q}$

Extending $f: H \to K$

Suffice to show one vertex extension, i.e.
 For ν being connected as φ_{ν,H} to w

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Claim provides a vertex w satisfying t over S and connected by 1 to F

Thus extending f

Thank you