

## Crossing Numbers: <br> Some Open Questions

Banff, 2018


## Crossing Numbers ?

grid crossing number, $21,26,44,71$
hierarchical crossing number, 68 hypercrossing, 70
hypergraph crossing numbers, 12, 23, 55
independent algebraic crossing number, 26 , $29,45,45,46$
independent crossing number, $4,4,8,11$, $22,26,45,69,71$
independent odd crossing number, $7,26,45$, 46, 57, 58, 71
independent odd projective plane crossing number, 46
independent pair crossing number, 26,46 , 59, 59
independent spherical crossing number, 4 independent string crossing number, 66 inner crossing number, 30 intersection-simple, 16, 40, 49, 50, 53
joint crossing numbers, $20,25,46$
-layer crossing number, $14,20,26,47,56$
$k$-page crossing number, 30
$k$-planar crossing number, 21, 27, 48
-quasi-planar, 38
Klein bottle crossing number, 21, 26, 35, 36 , 37
large angle crossing number, 63
eveled crossing number, $19,26,39,48,56$ 61
linear crossing number, $26,31,61$
local convex crossing number, 34,51
ocal crossing number, $5,11,17,22,26,49$ 57, 63
local outerplanar crossing number, 34, 35 local pair crossing number, 51 local toroidal crossing number, $11,26,49$
major crossing number, 26, 55 map crossing number, 21, 52, 70 maximal crossing number, 53 maximal rectilinear crossing number, 54
maximum convex rectilinear crossing number, 35
maximum crossing number, $7,8,22,26,53$, 54
maximum orchard crossing number, 22,26 , 58
maximum rectilinear crossing number, 5,7 , 22, 53, 53
maximum rectilinear edge crossing number, 26, 41
Metro-line crossing number, $15,23,33,54$, 70
minimum non-crossing edge number, 41
minor crossing number, $8,11,14,26,42,55$, 66
mixed upward crossing number, 68 monotone crossing number, $18,19,26,32$, $42,49,56,56,57,60$
monotone crossing numbers, 56,71
monotone independent odd crossing number, $11,26,57$
monotone odd crossing number, $5,17,26$, 56, 56, 57, 62
monotone odd + crossing number, 27
monotone odd $\pm$ crossing number, 27
monotone pair crossing number, $27,56,56$, 57, 59
monotone semisimple odd crossing number, 56
monotone weakly semisimple odd crossing number, 56
multiplanar crossing numbers, 49
nodal crossing number, $12,27,57$
nodal toroidal crossing number, 27, 57
non-crossing edge number, 41
non-orientable genus $g$ crossing number, 27
obfuscation complexity, 54
odd crossing number, $5,9,15,17,22,27$, $46,57,58,59,71$
odd + crossing number, $27,56,58,71$
odd $\pm$ crossing number, 56
odd $\pm$ crossing number, 58

## Solutions to Last Month's Puzzles.

 146.-WATER, GAS, AND ELECTRICITY.According to the conditions, in the strict sense in which one at first
 understands them, there is no possible solution to this puzzle. In such a dilemma one always has to look for some verbal quibble or trick. If the owner of house $\mathbf{A}$ will allow the water company to run their pipe for house $C$ through his property (and we are not bound to assume that he would object), then the difficulty is got over, as shown in our illustration. It will be seen that the dotted line from $W$ to $C$ passes through house A, but no pipe ever crosses another pipe.

## Definition

The (graph) crossing number, $\operatorname{cr}(G)$, of a graph $G$ is the smallest number of crossings in any drawing of $G$.

- drawn where? Plane, surface, book, ...
- drawn how?
- how is the graph represented?
- how is the representation visualized?
- how are crossings counted?


## Maximum Crossing Numbers

## Definition

## $\max -\overline{c r}(G)=$ largest \# crossings in

rectilinear (straight-line) drawing of $G$
$\max -\operatorname{cr}(G)=$ largest \# crossings in
good drawing of $G$

$\max -\operatorname{cr}\left(C_{4}\right)=\max -\overline{\operatorname{cr}}\left(C_{4}\right)=1$

## Maximum Crossing Numbers



$$
\max -\operatorname{cr}\left(K_{n}\right)=\max -\overline{c r}\left(K_{n}\right)=\binom{n}{4}
$$

## Polygons and Cycles



$$
\max -c r\left(C_{n}\right)=\frac{n(n-3)}{2} \text { for } n \geq 4
$$

$\max -\overline{\operatorname{cr}}\left(C_{n}\right)= \begin{cases}n(n-3) / 2 & \text { if } n \text { is odd }, \\ n(n-4) / 2+1 & \text { if } n \text { is even } .\end{cases}$

## Thrackle Bound

$$
\vartheta(G):=\frac{1}{2} \sum_{u v}(m-\operatorname{deg}(u)-\operatorname{deg}(v)+1)
$$

Lemma

$$
\max -\operatorname{cr}(G) \leq \vartheta(G)
$$

Theorem (Piazza, Ringeisen, Stueckle [PRS88] )

$$
\max -\operatorname{cr}(F)=\vartheta(F) \text { if } F \text { is a forest }
$$

Theorem (Verbitsky [V08])

$$
\frac{\vartheta(G)}{3} \leq \max -\overline{c r}^{o}(G) \leq \vartheta(G)
$$

## Maximum Crossing Number

Question: $\max -\overline{c r}^{o}(G)=\max -\overline{c r}(G)$ ?

Answer: no. [CFKUVW17]


What if G is bipartite? A tree? yes for spiders and trees of diameter at most 4 [BET18, FHKLSS18]

What if drawing is separated (all edges crossed by a line)?

Question: $\exists \mathbb{R}$-complete?
Known to be NP-hard [BJL16]

## Subgraph Monotonicity

Question: If $G$ is (induced) subgraph of $H$, then max-cr $(G) \leq$ max-cr $(H)$ ?

- True for max $\overline{c r}(G)$ (even pseudolinear $c r$ )
- Conditions on $G$ which make this happen?


## "Aura"/partial join Technique

## Theorem

If $C_{n}$ is induced subgraph of $H$, then $\max -c r\left(C_{n}\right) \leq \max -c r(H)$.


- True for odd $n$ and $C_{4}$
- True for $C_{6}$ (pic on left)
- Replace edge with $P_{3}$
- Add spokes


## "Aura"/partial join Technique

## Theorem

If $C_{n}$ is induced subgraph of $H$, then $\max -c r\left(C_{n}\right) \leq \max -c r(H)$.


Question: Is monotonicity conjecture true for (induced) subgraphs?

Exercise: Show that $\max -c r(G) \leq \max -c r(H)$ for $G$ apex, and induced subgraph of $H$.

## Maximum Crossing Number

Question: Is there a formula for $\max -\overline{C r}(T)$

$$
\text { if } T \text { is a tree? }
$$

Theorem (Piazza, Ringeisen, Stueckle [PRS88]) $\max -c r(F)=\vartheta(F)$ if $F$ is a forest

All caterpillars are thrackleable,
Subdivided $K_{1,3}$ is not.
yes for spiders and trees of diameter at most 4 [BET18, FHKLSS18]

## Conway’s Thrackle Conjecture

If a graph can be drawn in the plane so that every two edges intersect once, then $|E| \leq|V|$.

Current best bound: $|E| \leq 1.3984|V|$ (Fulek, Pach [FP17])

## Maximum Crossing Number Literature

[BJL16] Samuel Bald, Matthew P. Johnson, and Ou Liu. Approximating the maximum rectilinear crossing number. COCOON 2016, Springer, 2016.
[BET18] Patrick Bennett, Sean English, and Maria Talanda-Fisher. Weighted Turan Problems with Applications. ArXiv e-prints, abs/1809.05028, 2018. [CFKUVW17] Markus Chimani, Stefan Felsner, Stephen G. Kobourov, Torsten Ueckerdt, Pavel Valtr, and Alexander Wol?. On the maximum crossing number. CoRR, abs/1705.05176, 2017. arXiv:1705.05176
[FHKLSS18] Joshua Fallon, Kirsten Hogenson, Lauren Keough, Mario Lomelí, Marcus Schaefer, and Pablo Soberón. A Note on the Maximum Rectilinear Crossing Number of Spiders. ArXiv e-prints, abs/1808.00385, 2018.
[FP17] Radoslav Fulek, János Pach. Thrackles: An Improved Upper Bound. ArXiv e-prints, abs/2017arXiv170808037F, 2017.
[PRS88] B. L. Piazza, R. D. Ringeisen, and S. K. Stueckle. Properties of nonminimum crossings for some classes of graphs. In Graph theory, combinatorics, and applications. Vol. 2 (Kalamazoo, MI, 1988), Wiley-Intersci. Publ., pages 975 989 . Wiley, New York, 1991
[V08] Oleg Verbitsky. On the obfuscation complexity of planar graphs. Theoret. Comput. Sci., 396(1-3):294』300, 2008

## Local Crossing Number

## Definition

The local crossing number, lcr (D), of a drawing $D$ is the largest number of crossings along any edge in $D$.
The local crossing number, $\operatorname{lcr}(G)$, of a graph $G$ is the smallest $l c r(D)$ of anv drawing $D$ of $G$.


$$
\operatorname{lcr}\left(K_{6}\right)=1
$$


$\operatorname{lcr}\left(K_{7}\right)=2$

## Simple Local Crossing Number

## Definition

The simple local crossing number, $l c r^{*}(G)$, of a graph $G$ is the smallest $l \operatorname{cr}(D)$ of any good drawing $D$ of $G$.

Theorem (Pach, Radoičič, Tardos, Tóth [PRTT06]) If $\operatorname{lcr}(G) \leq 3$, then $\operatorname{lcr}(G)=l c r^{*}(G)$.

Used for better lower bound in crossing lemma

## Simple $=/ \neq$ Non-Simple

## Theorem

There is $G$ with $\operatorname{lcr}(G)=4$ and $l c r^{*}(G)=5$.


## Question: can $l c r^{*}(G)$ be bounded in $\operatorname{lcr}(G)$ ?

## Local Crossing Number



$$
\operatorname{lcr}\left(K_{6}\right)=1
$$

$$
\operatorname{lcr}\left(K_{7}\right)=2
$$

Question: $\operatorname{lcr}\left(K_{n}\right)$ ?

$$
\operatorname{lcr}\left(K_{m, n}\right) ?
$$

$\overline{l c r}\left(K_{n}\right)$ is known for all n (Ábrego, Fernandez-Merchant [AF17]) $\overline{\operatorname{lcr}}\left(K_{m, n}\right)$ is known for $\mathrm{m}=3,4$ [ADFLSS17].

## Local Crossing Number Literature

[AF17] Bernardo M. Ábrego and Silvia Fernández-Merchant. The rectilinear local crossing number of Kn. J. Combin. Theory Ser. A, 151:131?145, 2017.
[ADFLSS17] Bernardo M. Ábrego, Kory Dondzila, Silvia Fernández-Merchant, Evgeniya Lagoda, Seyed Sajjadi, and Yakov Sapozhnikov. On the rectilinear local crossing number of $K(m, n)$. Journal of Information Processing, 25:542?550, August 2017
[PRTT06] János Pach, Radoš Radoičić, Gábor Tardos, and Géza Tóth. Improving the crossing lemma by finding more crossings in sparse graphs. Discrete Comput. Geom., 36(4):527?552, 2006

## Independent Crossing Number

## Definition

The independent crossing number, $c_{-}(G)$, of a graph G is the smallest number of crossings between independent edges in any drawing of $G$.

Lemma $c r_{-}(G)=\operatorname{cr}(G)$
Proof


- fix drawing D of G with smallest number of crossings
- D cannot have dependent crossings
- so $c r_{-}(D)=\operatorname{cr}(D)$


## Independent Crossing Number

## Definition

The independent crossing number, $c r_{-}(G)$, of a graph G is ${ }^{\prime}$ '....' ${ }^{\text {I }}$ st number of crossings between independ Open $n$ any drawing of $G$.

Lemma $c r_{-}(G)=\operatorname{cr}(G)$

## Proof


fix drawingS of w with smallest number of

## Independent Crossing Number

Question: $\quad \operatorname{cr}(G)=c r_{-}(G)$ ?

- bound $c r$ in terms of $c r_{-}$(best bound quadratic)
- bound $\operatorname{cr}(D)$ of a $c r_{-}(G)$-minimal drawing (exponential bound possible)
- easier for $G=K_{n}$ ?


## Independent Crossing Number

## Conjecture:

If a graph can be drawn on a surface without independent crossings. Then graph can be embedded in surface.

Known for plane (Hanani-Tutte theorem) projective plane (HT for PP)
HT fails for orientable surfaces of genus $\geq 4$.

## Rectilinear Crossing number

## Definition

The rectilinear crossing number, $\overline{c r}(G)$, of a graph G is the smallest number of crossings in any straight-line drawing of G .

Theorem (Fary, 48; Wagner, 36)
If $G$ is planar, then $G$ has a planar straight-line drawing.

Theorem (Bienstock, Dean, 93)

$$
\overline{c r}(G)=\operatorname{cr}(G) \text { for } \operatorname{cr}(G) \leq 3
$$

## Independent Crossing Number Literature

[S10] Marcus Schaefer. Removing Incident Crosssings. Manuscript, 2010.

## Rectilinear Crossing number

## Definition

The rectilinear crossing number, $\overline{c r}(G)$, of a graph G is the smallest number of crossings in any straight-line drawing of G .

Theorem (Fary, 48; Wagner, 36)
If $G$ is planar, then $G$ has a planar straight-line drawing.

Theorem (Bienstock, Dean, 93)

$$
\overline{c r}(G)=\operatorname{cr}(G) \text { for } \operatorname{cr}(G) \leq 3
$$

## Conjectures and Question

Conjecture (Harary, Kainen, Schwenk [HLS73]) $\overline{c r}\left(C_{m} \square C_{n}\right)=n(m-2)$ for $n \geq m \geq 3$.

Also open for $\operatorname{cr}(G)$. For that case, partial results known.

Conjecture (Hernández-Vélez, Leaños, Salazar [HLS17] ) $\overline{c r}(G)$ can be bounded in $\operatorname{cr}(G)$ for 3-connected $G$.

Question: Can $\overline{c r}$ be bounded in $\widetilde{c r}(G)$ ?
Question: What's the complexity of $\overline{c r}(G) \leq 4$ ?

## Grid Drawings

Question: Is there an $f$ so that
if $\overline{c r}(G) \leq k$ and $n=|V(G)|$, then
$G$ can be realized on a $f(k) n \times f(k) n$ grid?

Alternatively: on an $n^{f(k)} \times n^{f(k)}$ grid?

Question: If we restrict drawing $D$ of $G$ to a $t \times t$ grid, what is the best $\overline{c r}(D)$ we can guarantee?

$$
\text { If } t=\Omega\left(2^{2^{n}}\right) \text {, then we can achieve } \overline{c r}(G)
$$

## Rectilinear Crossing Number Literature

[BD93] Daniel Bienstock and Nathaniel Dean. Bounds for rectilinear crossing numbers. J. Graph Theory, 17(3):333国348, 1993.
[HKS73] Frank Harary, Paul C. Kainen, and Allen J. Schwenk. Toroidal graphs with arbitrarily high crossing numbers. Nanta Math., 6(1):58@67, 1973.
[HLS17] César Hernández-Vélez, Jesús Leaños, and Gelasio Salazar. On the pseudolinear crossing number. Journal of Graph Theory, 84(3):297国310, 2017.

## Counting Crossings

$\operatorname{cr}(D)=\#$ of crossings in $D$

$\operatorname{pcr}(D)=\# r$ of pairs of edges crossing in $D$
$\operatorname{ocr}(D)=\#$ of pairs of edges crossing oddly in $D$
ocr_ $(D)=\#$ of pairs of independent edges crossing oddly in $D$

$$
\operatorname{ocr}_{-}(G) \leq \operatorname{ocr}(G) \leq \operatorname{pcr}(G) \leq \operatorname{cr}(G)
$$

## The Hanani-Tutte Theorem

Theorem (Hanani, 1934; Tutte, 1970)

$$
\operatorname{ocr}_{-}(G)=0 \text { implies } \operatorname{cr}(G)=0
$$



## $\operatorname{ocr}_{-}(G)=\operatorname{ocr}(G)=\operatorname{pcr}(G)=\operatorname{cr}(G) ?$

Answer: Not all. [FPSS12]
There is $G$ so that $\operatorname{ocr}_{-}(G)<\operatorname{ocr}(G)<\operatorname{pcr}(G)$

Conjecture $\operatorname{pcr}(G)=\operatorname{cr}(G)$

## Bounds?

Theorem (Matousek [M14])

$$
\begin{aligned}
& \operatorname{cr}(G) \leq \operatorname{pcr}(G)^{\frac{3}{2}} \log ^{2}(\operatorname{pcr}(G)) \\
& \quad \log ^{2}(\operatorname{pcr}(G)) \text { can be improved to } \log (\operatorname{pcr}(G)) \\
& \quad \text { using string graph separator by Lee. }
\end{aligned}
$$

Theorem (Pelsmajer, Schaefer, Štefankovič [PSS10])

$$
\operatorname{cr}(G) \leq\binom{ 2 o c r_{-}(G)}{2}
$$

Question: sub-quadratic bounds for $\mathbf{c r}$ in ocr_?

## Crossing Lemma for $\operatorname{pcr}(G)$

remove each vertex of $G=(V, E)$ with prob. $p$ i.a.r.

$$
\begin{aligned}
& G^{\prime}=\left(V^{\prime}, E^{\prime}\right) \quad \operatorname{cr}\left(G^{\prime}\right) \geq\left|E^{\prime}\right|-3\left|V^{\prime}\right| \\
& \mathbb{E}\left(\operatorname{cr}\left(G^{\prime}\right)\right) \geq \mathbb{E}\left(\left|E^{\prime}\right|\right)-3 \mathbb{E}\left(\left|V^{\prime}\right|\right) \\
& p^{4} \operatorname{cr}(G) \geq p^{2}|E|-3 p|V|
\end{aligned}
$$

can replace cr with pcr?

$$
\begin{aligned}
& \operatorname{cr}(G) \geq p^{-2}|E|-3 p^{-3}|V| \\
& \operatorname{cr}(G) \geq \frac{|E|^{3}}{64|V|^{2}} \quad \text { e.g. } p=4|V| /|E|
\end{aligned}
$$

## Improved Crossing Lemma for $\operatorname{pcr}_{+}(G)$

 $p c r_{+}$: no adjacent crossings allowedTheorem (Ackerman, Schaefer [AS14])

$$
\operatorname{pcr}_{+}(G) \geq \frac{1}{32.4} \frac{m^{3}}{n^{2}} \text { for } m \geq 6.75 n
$$

use $\operatorname{lpcr}(G) \leq 2$, then $\operatorname{lcr}(G) \leq \operatorname{lpcr}(G)$

Questions:

$$
\begin{aligned}
& \operatorname{lpcr}(G) \leq 3, \text { then } \operatorname{lcr}(G) \leq \operatorname{lpcr}(G) \\
& \operatorname{lpcr}_{-}(G) \leq 1, \text { then } \operatorname{lcr}(G) \leq 1 ?
\end{aligned}
$$

## Odd and Pair Crossing Number Literature

[AS14] Eyal Ackerman and Marcus Schaefer. A crossing lemma for the pair-crossing number. Graph Drawing, 222-233, 2014.
[FPSS12] Radoslav Fulek, Michael J. Pelsmajer, Marcus Schaefer, and Daniel Štefankovič, Adjacent crossings do matter. J. Graph Algorithms Appl., 16(3):759?782, 2012.
[PSS10] Michael Pelsmajer, Marcus Schaefer, Daniel Štefankovič. Removing independently even crossings. SIAM Journal on Discrete Mathematics, 24(2):379@393, 2010

## Thank You

## Open Questions on Crossing Numbers

[A09] Dan Archdeacon. Open problems. In Topics in topological graph theory, volume 128 of Encyclopedia Math. Appl., Cambridge, 2009.
[BMP05] Peter Brass, William Moser, and János Pach. Research Problems in Discrete Geometry. Springer, New York, 2005.
[PTO0] János Pach and Géza Tóth. Thirteen problems on crossing numbers. Geombinatorics, 9(4):194?207, 2000.
[S17] Marcus Schaefer. The Graph Crossing Number and its Variants: A Survey, Electronic Journal of Combinatorics, 2017.
[S16] László A. Székely. Turán's brick factory problem: The status of the conjectures of
Zarankiewicz and Hill. In Ralucca Gera, Stephen Hedetniemi, and Craig Larson, editors, Graph Theory: Favorite Conjectures and Open Problems, Springer, 2016.

Also:

- http://www.openproblemgarden.org/category/topological graph theory
- http://www.cems.uvm.edu/TopologicalGraphTheoryProblems/



## Question: What is the oldest reference to crossing numbers?

From de la Vera Cruz' Recognitio Summularum, 1554 (http://www.primeroslibros.org/ browse.html)


From Mary L. Northway. A Method for Depicting Social Relationships Obtained by Sociometric Testing. Sociometry, Vol. 3, No. 2 (Apr., 1940), pp. 144-150

