

# Relation Generation in Quadratic Number and Function Fields

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# Imaginary Quadratic Number Fields

$\mathbb{Q}(\sqrt{\Delta}) = \{x + y\sqrt{\Delta} \mid x, y \in \mathbb{Q}\}$  : quadratic field

- $\Delta \equiv 0, 1 \pmod{4}$  : discriminant ( $\in \mathbb{Z}$ ,  $\Delta$  or  $\Delta/4$  square-free)
- $\Delta < 0$  : *imaginary* quadratic field

$\mathcal{O}_\Delta \subset \mathbb{Q}(\sqrt{\Delta})$  : maximal order of  $\mathbb{Q}(\sqrt{\Delta})$  (ring of algebraic integers)

- $\mathcal{I}_\Delta$  : group of invertible, fractional ideals of  $\mathcal{O}_\Delta$
- $\mathcal{P}_\Delta$  : principal, fractional ideals, subgroup of  $\mathcal{I}_\Delta$
- $Cl_\Delta = \mathcal{I}_\Delta / \mathcal{P}_\Delta$  : class group
- $h_\Delta = |Cl_\Delta|$  : class number
- unique reduced ideal representatives of group elements

# Relations

*Relation*: power-product of prime ideals that is principal

Used in index-calculus algorithms for:

- invariant computation (class number, class group structure, regulator/fundamental unit)
- discrete logarithm computation, principality testing / norm equations
- computing large-degree isogenies and endomorphism rings of ordinary elliptic curves over finite fields

Efficiency of all depends on quickly finding relations

# Example: Computing the Class Group

## Outline:

- factor base  $FB$  : prime ideals  $\mathfrak{p}_i$  of norm  $p_i \leq B$ , must generate  $Cl_\Delta$
- surjective homomorphism (assume  $|FB| = k$ )

$$\begin{aligned} \varphi : \mathbb{Z}^k &\rightarrow Cl_\Delta \\ (v_1, \dots, v_k) &\mapsto [\mathfrak{p}_1^{v_1} \dots \mathfrak{p}_k^{v_k}] \end{aligned}$$

- $\mathbb{Z}^k / \Lambda \cong Cl_\Delta$ , where  $\Lambda = \ker \varphi$  is the lattice of all relations wrt  $FB$
- randomly construct generating system of  $\Lambda$ , linear algebra (Smith normal form) to compute group structure

Expected run time (GHR):  $L_\Delta(1/2, \sqrt{2})$ , where

$$L_\Delta(\alpha, \beta) = \exp((\beta + o(1))(\log|\Delta|)^\alpha (\log \log|\Delta|)^{1-\alpha})$$

# Example: Computing Large-Degree Isogenies

$Ell_{t,u}(\mathbb{F}_q)$  : isomorphism classes of elliptic curves over  $\mathbb{F}_q$  with trace  $t$  and endomorphism ring  $\mathcal{O}_{u^2\Delta_K} \in \mathbb{Q}(\sqrt{\Delta_K})$

## Theorem

*Let  $\mathfrak{a} \subset \mathcal{O}_{u^2\Delta_K}$  be prime of norm  $\ell$ . Then  $\mathfrak{a}$  acts on  $Ell_{t,u}(\mathbb{F}_q)$  via a degree  $\ell$  isogeny, defining a faithful group action by  $Cl_{u^2\Delta_K}$ .*

Jao, Soukharev 2010: idea (compute isogeny of degree  $\ell$ ):

- Compute relation  $\mathfrak{p}_\ell \prod \mathfrak{p}_i^{e_i}$  in  $Cl_{u^2\Delta_K}$  for  $\mathfrak{p}_i$  small,  $N(\mathfrak{p}_\ell) = \ell$
- $[\mathfrak{p}_\ell] = \prod [\mathfrak{p}_i]^{-e_i} \in Cl_{u^2\Delta_K}$
- Evaluate the degree  $\ell$  isogeny via evaluations of degree  $\mathfrak{p}_i$  isogenies

Expected run time (GRH):  $L_q(1/2, \sqrt{3}/2) \log \ell$

# Finding Relations

Main idea:

- Compute  $\mathfrak{a} \sim \prod p_i^{e_i}$  (but not equal!)
- If  $\mathfrak{a} = \prod p_i^{v_i}$ , then  $\prod p_i^{e_i - v_i}$  is principal

One approach: random selection of  $\mathfrak{a}$  via choice of  $e_i$  (or random walks)

Better approach: sieving

- let  $\alpha = ax + (b + \sqrt{\Delta})/2y \in \mathfrak{a} = a\mathbb{Z} + (b + \sqrt{\Delta})/2\mathbb{Z}$
- $N(\alpha) = a(ax^2 + bxy + cy^2)$  where  $c = (b^2 - \Delta)/(4a)$
- there exists ideal  $\mathfrak{b}$  with  $N(\mathfrak{b}) = ax^2 + bxy + cy^2$  and  $(\alpha) = \mathfrak{a}\mathfrak{b}$
- find  $x, y \in \mathbb{Z}$  such that  $f(x, y) = ax^2 + bxy + cy^2$  factors over the  $p_i$

# Sieving

Finding relations  $\leftrightarrow$  finding smooth values of  $f(X, Y) = aX^2 + bXY + cY^2$

One approach: find all  $x \leq M$ ,  $x \in \mathbb{Z}$ , with  $f(x, 1) = ax^2 + bx + c$  smooth

For each prime ideal of norm  $p_i$  :

- compute root(s)  $r$  such that  $f(r, 1) \equiv 0 \pmod{p_i}$
- $p_i \mid r$ , and  $p \mid kp_i + r$  for all  $k \in \mathbb{Z}$
- use analogue of Sieve of Eratosthenes to factor all  $f(x, 1)$  by “marking off” every  $p_i$ th cell in an array, starting at  $r$

Can adapt quadratic sieve methods from integer factoring, including self-initialization

# Some Results

Biasse (2010): class group for  $\Delta = -4 \times 10^{110} - 4$

$$Cl_{\Delta} \cong \mathbb{Z}/8576403641950292891121955131452148838284294200071440\mathbb{Z} \times (\mathbb{Z}/2\mathbb{Z})^{11}$$

Biasse, J. (2010): class group and regulator for  $\Delta = 4 \times 10^{110} + 4$

$$Cl_{\Delta} \cong \mathbb{Z}/12\mathbb{Z} \times \mathbb{Z}/2\mathbb{Z}$$

$$R_{\Delta} \approx 70795074091059722608293227655184666748799878533480399.67302$$

4 days for relations (260 2.4 GHz Xeons), 4 days for linear algebra (2.4 GHz Opteron, 32 GB RAM), 4 days for GRH-verification



# Isogeny and Endomorphism Ring Computation: Obstacles

Parameter tuning is really hard

- Composition of factor base can affect results dramatically
- Eg. (J. 1999), computing  $Cl_{\Delta}$ 
  - typical 70-decimal digit  $\Delta$  :  $18h$
  - 70-decimal digit  $\Delta$  with no  $p_i \leq 353$  in factor base: 6.5 days

Need really small factor bases for isogeny and endomorphism ring computation

- only small prime degree isogenies are efficient to compute
- sieving becomes more effective with larger factor bases

# Our Approach (on-going work)

Analytic model to estimate smoothness probabilities given a particular factor base

- extend numerical methods to approximate  $\psi(x, y)$  to ideals of quadratic fields
- would take into account differing splitting behavior of small primes
- use as basis of search for optimal parameters

Use Sutherland's improvements to evaluation of low-degree isogenies

- feasible to evaluate isogenies of larger prime degree
- may be sufficient to realize benefits from sieving

# Imaginary Quadratic Function Fields

$C : y^2 + h(t)y = f(t)$  non-singular,  $h, f \in \mathbb{F}_q[t]$

$C$  is *imaginary* (genus  $g$ ) if

- $q$  is odd,  $h = 0$ ,  $f$  monic and square-free with  $\deg(f) = 2g + 1$
- $q$  is even,  $h \neq 0$  with  $\deg(h) \leq g$  and  $f$  monic with  $\deg(f) = 2g + 1$

(a.k.a. hyperelliptic curves)

$\deg 0$  divisor class group (ideal class group of  $\mathbb{F}_q(C)$ ):

- finite abelian, size  $\approx q^g$
- unique reduced divisor/ideal representatives of group elements

# Example Application: Weil Descent

Reduce elliptic curve discrete logarithm problem (over  $\mathbb{F}_{2^{ng}}$ ) to hyperelliptic curve discrete logarithm problem (genus  $g$  over  $\mathbb{F}_{2^n}$ )

- Enge, Gaudry (index-calculus): if  $g > \log q$ , expected run time  $L_{q^g}(1/2, 5.73 + o(1))$
- J, Menezes, Stein: implementation, parameter optimization
  - solved ECDLP over  $\mathbb{F}_{2^{31}}$ ,  $\mathbb{F}_{2^{64}}$ ,  $\mathbb{F}_{2^{93}}$ , and  $\mathbb{F}_{2^{124}}$
  - genus 31 hyperelliptic curves defined over  $\mathbb{F}_2$ ,  $\mathbb{F}_{2^2}$ ,  $\mathbb{F}_{2^3}$ , and  $\mathbb{F}_{2^4}$
- Velichka, J., Stein: application of sieving, solved ECDLP over  $\mathbb{F}_{2^{155}}$ 
  - genus 31 hyperelliptic curve defined over  $\mathbb{F}_{2^5}$

# Overview of Index Calculus and Sieving

Same general approach as in quadratic fields

- factor base: prime ideals  $\mathfrak{p}$  with  $\deg \mathfrak{p}_i \leq B$  ( $p_i$  irreducible)
- find random relations
- solve linear algebra problem (linear system modulo group order)

Can apply same approach to finding relations, including sieving

- relation generation reduces to finding smooth values of  $f(X) = aX^2 + bX + c$  defined over  $\mathbb{F}_q[t]$
- same improvements (eg. self-initialization) are possible

# Challenges with Sieving

Need to find all  $x \in \mathbb{F}_q[t]$  with  $\deg(x) \leq M$  such that  $f(x)$  is  $B$ -smooth

How to map  $x \in \mathbb{F}_q[t]$  to a cell in an array?

- Natural map (Flassenburg, Paulus 1998),  $q = p^d$  :

$$\nu : \mathbb{F}_q[t] \rightarrow \mathbb{Z}$$

$$x_m t^m + \cdots + x_0 \mapsto \nu_0(x_m) q^m + \cdots + \nu_0(x_0)$$

where

$$\nu_0 : \mathbb{F}_q \rightarrow \{0, \dots, q-1\}$$

$$\nu_0(a_d \alpha^d + \cdots + a_0) = a_d p^d + \cdots + a_0$$

Works, but painful to evaluate frequently

# Challenges with Sieving, cont.

For irreducible  $p_i \in \mathbb{F}_q[t]$  and  $r \in \mathbb{F}_q[t]$  such that  $f(r) \equiv 0 \pmod{p_i}$  :

- how to rapidly find all  $\nu(kp_i + r)$  for  $k \in \mathbb{F}_q[t]$  such that  $\deg(kp_i + r) \leq M$ ?
- map  $\nu$  does not lead to regular spacing through the sieve array

Velichka, J., Stein 2008: enumerate all  $k$  of appropriate degree, evaluate  $\nu(kp_i + r)$  directly using previous results and precomputations

- use  $k'p_i + r = (kp_i + r) + (k' - k)p_i$  (add appropriate multiple of  $p$ )

Trei, J. Stein 2013: further optimizations, including

- evaluation at  $q$  using Horner's rule
- better use of intermediate results
- observation that  $\nu(x + y) = \nu(x) \oplus \nu(y)$  (all ops on integers)

# Numerical Results

VJS 2008 results (278 Intel P4 Xeon 2.4 GHz CPUs, 26 2.8 GHz):

- ECDLP over  $\mathbb{F}_{2^{124}}$  (HCDLP with  $g = 31$ ,  $q = 2^4$ ):
  - 9 hours, 7.5 hours for relations (24 hours with random walks)
- First solution of ECDLP over  $\mathbb{F}_{2^{155}}$  (HCDLP with  $g = 31$ ,  $q = 2^5$ ):
  - 3 weeks, 1 week for relations (random walks estimate 5 weeks)

TJS 2013 results (64 Intel Xeon X7560 2.27 GHz CPUs):

- $\mathbb{F}_{2^{124}}$  : 3 hours (27 min. for relations)
- $\mathbb{F}_{2^{155}}$  : in progress (2.5 days for relations)



# Future Work

Complete analytic model to aid parameter selection

Two dimensional (lattice) sieving?

Batch smoothness test for candidates produced by the sieve?

Function fields:

- add double large primes
- try odd characteristic
- lower genus?