Jonathan Hauenstein Convex Algebraic Geometry BIRS February 16, 2010

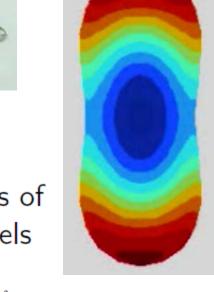




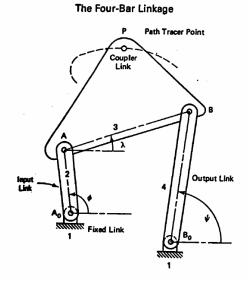


Inverse 6R problem

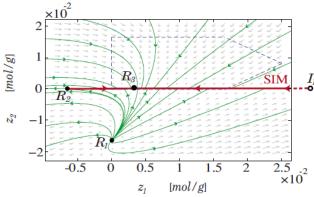
Zebrafish patterning



Bifurcations of tumor models

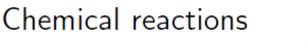


Nine-point problem



 $H_2+O_2 \rightleftharpoons 2OH$ $OH+H_2 \rightleftharpoons H_2O+H$ $H+O_2 \rightleftharpoons OH+O$ $O+H_2 \rightleftharpoons OH+H$ $H+O_2+M \rightleftharpoons HO_2+M^b$ $H+2O_2 \rightleftharpoons HO_2+O_2$ $H+O_2+N_2 \rightleftharpoons HO_2+N_2$ $OH+HO_2 \rightleftharpoons H_2O+O_2$ $H+HO_2 \rightleftharpoons 2OH$ $O+HO_2 \rightleftharpoons O_2+OH$ $2OH \rightleftharpoons O + H_2O$ $H_2+M \rightleftharpoons 2H+M^c$ $O_2+M \rightleftharpoons 2O+M$ $H+OH+M \rightleftharpoons H_2O+M^d$ $H+HO_2 \rightleftharpoons H_2+O_2$ $2HO_2 \rightleftharpoons H_2O_2 + O_2$ $H_2O_2+M \rightleftharpoons 2OH+M$ $H_2O_2+H\rightleftharpoons HO_2+H_2$ $H_2O_2+OH \rightleftharpoons H_2O+HO_2$





Overview

- Homotopy continuation
- Basic numerical algebraic geometry
- Regeneration
- Rank-deficiency sets

Joint work with

- D. Bates (Colorado State)
- C. Peterson (Colorado State)
- A. Sommese (Notre Dame)
- C. Wampler (General Motors R&D)





General references

The Numerical Solution of Systems of Polynomials

Arising in Engineering and Science









Andrew J. Sommese - Charles W. Wampler, II

T.Y. Li, Numerical solution of polynomial systems by homotopy continuation methods, in *Handbook of Numerical Analysis*, Volume XI, 209–304, North-Holland, 2003.



General references

D.J. Bates, J.D. Hauenstein, A.J. Sommese, and C.W. Wampler, Bertini: Software for Numerical Algebraic Geometry. Available at www.nd.edu/~sommese/bertini.













Main problem in numerical algebraic geometry: Describe all $x \in \mathbb{C}^N$ where

$$f(x) = \begin{bmatrix} f_1(x_1, \dots, x_N) \\ \vdots \\ f_n(x_1, \dots, x_N) \end{bmatrix} = 0$$

and each f_i is polynomial.





Basic isolated root finding:

Assume n = N ("square"). Compute the isolated solutions of

$$f(x) = \begin{bmatrix} f_1(x_1, \dots, x_N) \\ \vdots \\ f_n(x_1, \dots, x_N) \end{bmatrix} = 0.$$





Algorithm

- ▶ Treat f as a member of a parameterized family of polynomial systems \mathcal{F} .
- ▶ Compute the isolated roots of $g \in \mathcal{F}$ (general enough).
- ▶ Setup the homotopy H(x, t) = (1 t)f(x) + tg(x).
- ► Track the paths x(t) defined by $H(x(t), t) \equiv 0$. Since H(x(1), 1) = g(x(1)) = 0, paths start at the known roots of g.





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Computing the roots of g is very interesting when nontrivial.

B. Huber and B. Sturmfels, A polyhedral method for solving sparse polynomial systems, *Math. Comp.* 64(212), 1541–1555, 1995.





$$f = \begin{bmatrix} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{bmatrix}$$

1. Total degree:
$$\mathcal{F} = \left\{ \begin{bmatrix} g_1(x,y) \\ g_2(x,y) \end{bmatrix} : \deg(g_i) = 2 \right\}$$
,

$$g = \begin{bmatrix} x^2 - 1 \\ y^2 - 1 \end{bmatrix}$$
, Bound: 4.





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, Bound: 4.

2. 2-hom:
$$\mathcal{F} = \left\{ \begin{bmatrix} g_1(x) \\ g_2(x,y) \end{bmatrix} : \deg(g_1) = 2 \\ \deg_x(g_2) = \deg_y(g_2) = 1 \right\},$$

$$g = \begin{bmatrix} x^2 - 1 \\ (x - 2)(y - 1) \end{bmatrix}, \text{ Bound: 2}.$$





$$f = \begin{bmatrix} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{bmatrix}$$

3. Polytope:
$$\mathcal{F} = \left\{ \begin{bmatrix} a_1x^2 + a_2x + a_3 \\ a_4xy + a_5x + a_6y + a_7 \end{bmatrix} : a_i \in \mathbb{C} \right\}$$
, $g = \begin{bmatrix} x^2 - 1 \\ y - 1 \end{bmatrix}$, Bound: 2.





$$f = \begin{bmatrix} x^2 + 2x - 8 \\ xy + 2x + 4y - 3 \end{bmatrix}$$

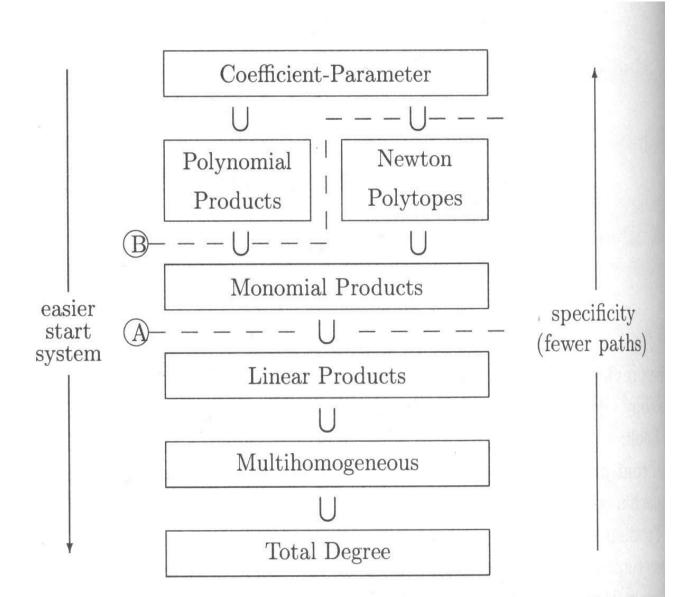
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$$\mathcal{F} = \left\{ \begin{bmatrix} a_1x^2 + a_2x + a_3 \\ a_4xy + a_5x + a_6y + a_7 \end{bmatrix} : a_i \in \mathbb{C} \right\}$$
, $g = \begin{bmatrix} x^2 - 1 \\ y - 1 \end{bmatrix}$, Bound: 2.

4. "Optimal":
$$\mathcal{F} = \left\{ \begin{bmatrix} x^2 - (a_1 + a_2)x + a_1a_2 \\ (x - a_1)y + a_3x + a_4 \end{bmatrix} : a_i \in \mathbb{C} \right\},$$

$$g = \begin{bmatrix} x^2 - 1 \\ (x + 1)y - 2 \end{bmatrix}, \text{ Bound: } 1.$$

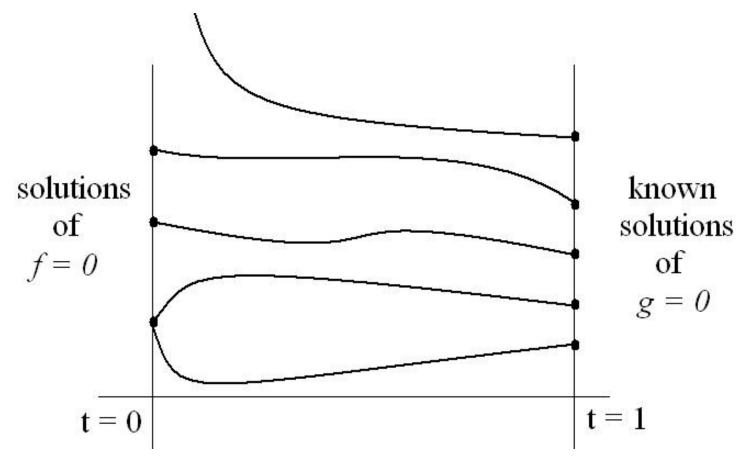












$$H(x,t) = (1-t)f(x) + tg(x) \equiv 0$$





For a properly constructed homotopy:

- ▶ Solution paths x(t) exist
- \triangleright Solution paths x(t) satisfy the Davidenko differential equation

$$0 \equiv \frac{dH(x(t),t)}{dt} = \frac{\partial H(x(t),t)}{\partial x}x'(t) + \frac{\partial H(x(t),t)}{\partial t}.$$

- ▶ For $t \neq 0$, $\frac{\partial H(x(t),t)}{\partial x}$ is invertible.
- ▶ {isolated roots of f} \subset $\{x(0) = \lim_{t\to 0} x(t) \mid x(1) \text{ is an isolated root of } g\}$





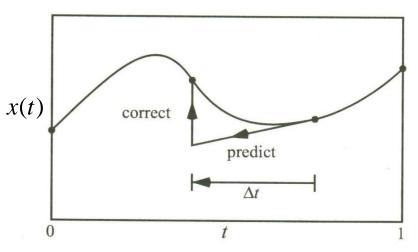
Track solution paths using a predictor-corrector scheme.

Predict using Davidenko's differential equation:

$$\frac{\partial H(x(t),t)}{\partial x}x'(t) = -\frac{\partial H(x(t),t)}{\partial t}.$$

Correct using Newton's method:

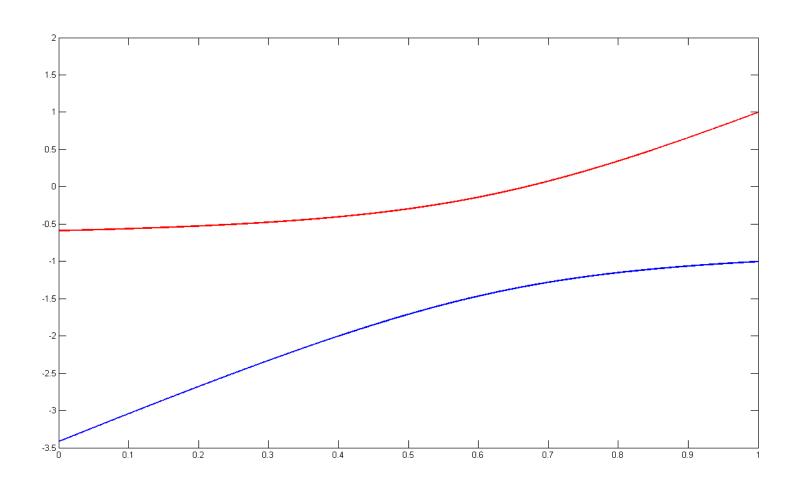
$$H(x(t), t) \equiv 0.$$





$$f(x) = x^2 + 4x + 2$$

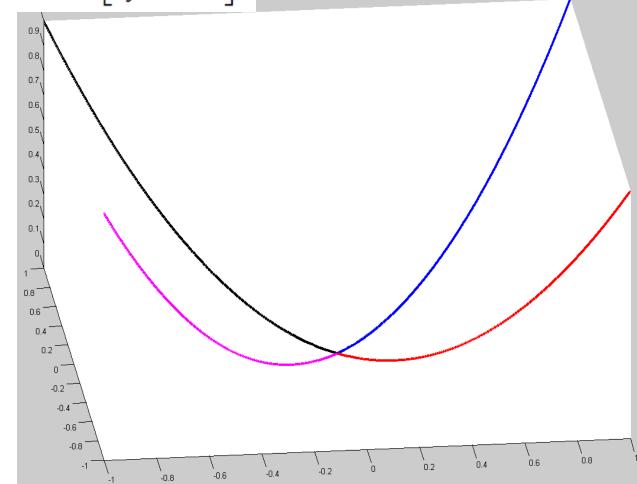
$$H(x, t) = (1 - t)f(x) + t(x^2 - 1).$$





Example
$$f(x,y) = \begin{bmatrix} x^2 \\ y^2 \end{bmatrix}$$

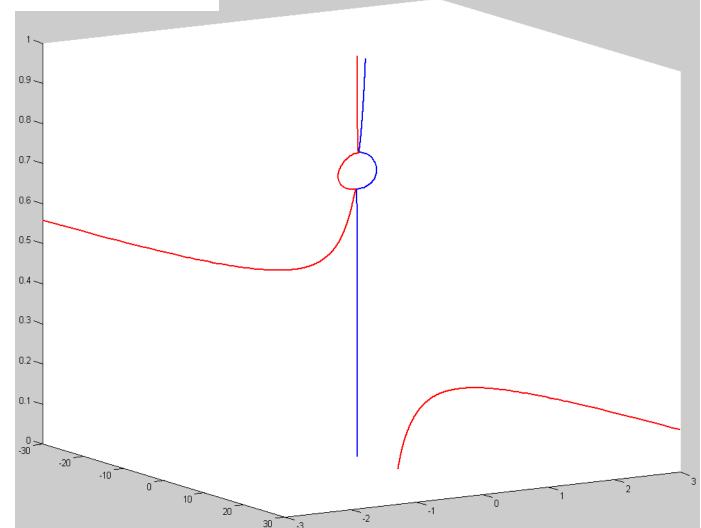
$$H(x, y, t) = (1 - t)f(x, y) + t \begin{bmatrix} x^2 - 1 \\ y^2 - 1 \end{bmatrix}.$$





$$f(x) = -\frac{1}{2}x^2 + 4x + \frac{14}{3}$$

$$H(x, t) = (1 - t)f(x) + t(x^2 - 1).$$



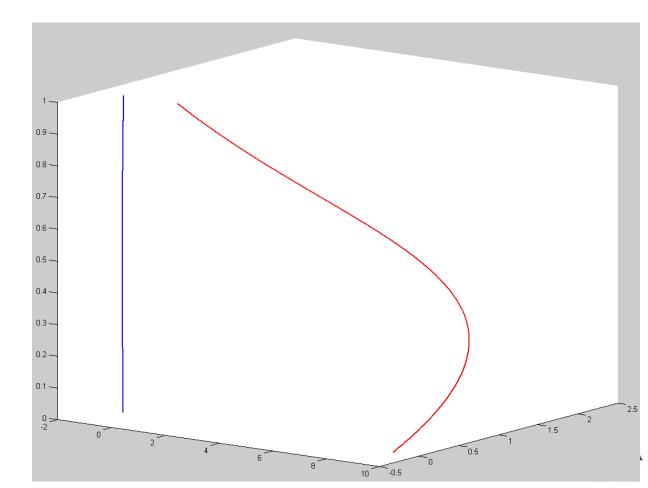


Example

$$f(x) = -\frac{1}{2}x^2 + 4x + \frac{14}{3}$$

For random $\gamma \in \mathbb{C}$,

$$H(x, t) = (1 - t)f(x) + \gamma t(x^2 - 1).$$





Singular endpoints occur frequently.

- ▶ Endgames: compute the endpoint by staying sufficiently far away from t = 0.
- ▶ Deflation: restore quadratic convergence of Newton iterations.

$$f = \begin{bmatrix} x^2 \\ y^2 \end{bmatrix} H = \begin{bmatrix} x^2 - t \\ y^2 - t \end{bmatrix}$$
$$J = \begin{bmatrix} 2x & 0 \\ 0 & 2y \end{bmatrix}.$$

$$J \rightarrow 0$$
 as $t \rightarrow 0$.

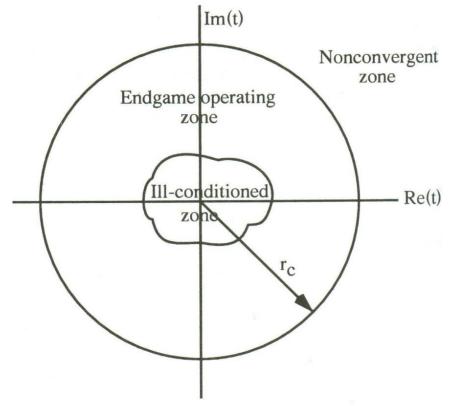




Endgame algorithms accurately compute the endpoint of the path by using the local Puiseux series expansion:

$$x(t)=x(0)+\sum_{j\geq 1}a_jt^{j/c}.$$

Use high enough precision to ensure reliable numerical computations.



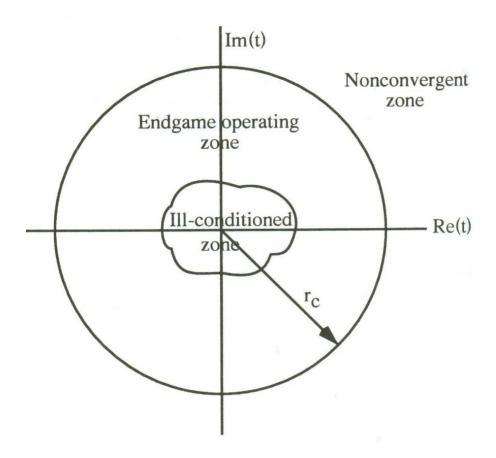




$$x(t) = x(0) + \sum_{j \geq 1} a_j t^{j/c}$$

Cauchy integral theorem:

$$x(0) = \frac{1}{2\pi c} \int_0^{2\pi c} x(Re^{i\theta}) d\theta.$$







Deflation for isolated solutions

A. Leykin, J. Verschelde, and A. Zhao, Newton's method with deflation for isolated singularities of polynomial systems, *Theor. Comput. Sci.*, 359, 111–122, 2006.

$$\begin{bmatrix} x^{2} \\ y^{2} \end{bmatrix} \implies \begin{bmatrix} x^{2} \\ y^{2} \\ 2x\lambda_{1} \\ 2y\lambda_{2} \\ \alpha_{1}\lambda_{1} + \alpha_{2}\lambda_{2} - 1 \\ \beta_{1}\lambda_{1} + \beta_{2}\lambda_{2} - 1 \end{bmatrix}$$

$$(0,0), \text{ mult 4}$$

$$\begin{pmatrix} 0,0,\widehat{\lambda_{1}},\widehat{\lambda_{2}} \end{pmatrix}, \text{ mult 1}$$





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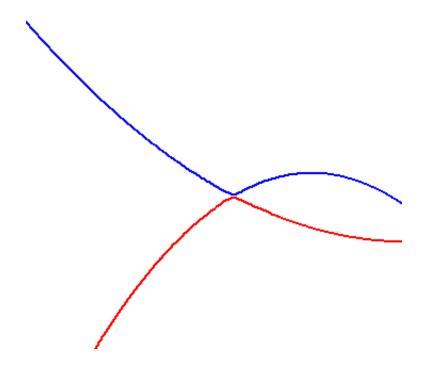
- ▶ For $t \neq 0$, $\frac{\partial H(x(t),t)}{\partial x}$ is invertible.
 - ▶ {isolated roots of f} \subset $\{x(0) = \lim_{t\to 0} x(t) \mid x(1) \text{ is an isolated root of } g\}$





Near singularities arise often that can make path tracking numerically challenging.

- D.J. Bates, J.D. Hauenstein, A.J. Sommese, and C.W. Wampler, Adaptive multiprecision path tracking. *SIAM J. Num. Anal.*, 46(2), 722–746, 2008.
- D.J. Bates, J.D. Hauenstein, A.J. Sommese, and C.W. Wampler, Stepsize control for adaptive multiprecision path tracking. *Contemp. Math.*, 496, 21–31, 2009.





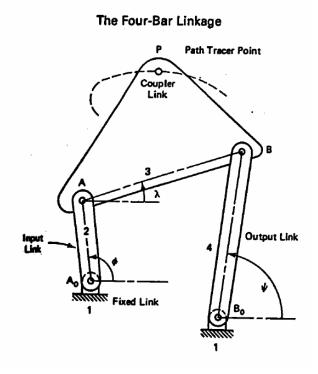


Applied the method of

D.J. Bates, J.D. Hauenstein, A.J. Sommese, and C.W. Wampler, Stepsize control for adaptive multiprecision path tracking. *Contemp. Math.*, 496, 21–31, 2009.

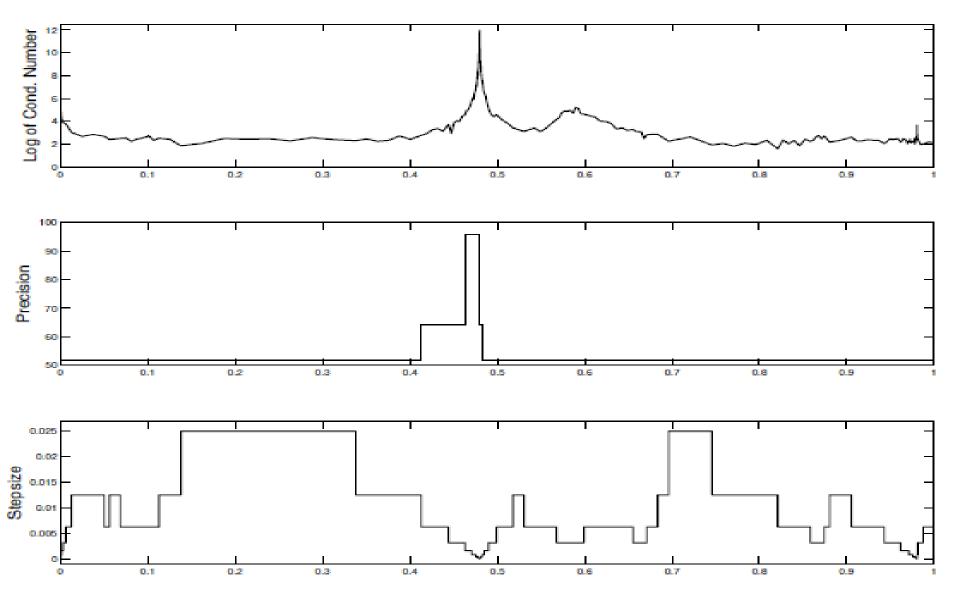
to the nine-point problem for four-bar linkages.

Out of 143,360 paths, 1184 paths (0.83%) needed higher precision to successfully track past near-singularity conditions.









D.J. Bates, J.D. Hauenstein, A.J. Sommese, and C.W. Wampler, Stepsize control for adaptive multiprecision path tracking. *Contemp. Math.*, 496, 21–31, 2009.





Using only double precision is fast, but can lead to path crossings

T.L. Lee, T.Y. Li, C.H. Tsai, HOM4PS-2.0: a software package for solving polynomial systems by the polyhedral homotopy continuation method. *Computing*, 83(2-3), 109–133, 2008.

and results may not be correct.

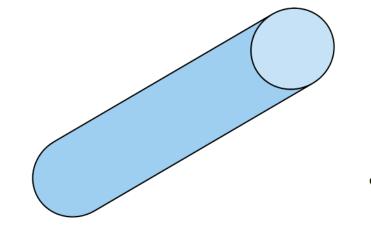
H. Tari, H.J. Su, and T.Y. Li, A constrained homotopy technique for excluding unwanted solutions from polynomial equations arising in kinematics problems. To appear in *Mechanism and Machine Theory*.





Positive-dimensional solution sets can be handled by intersecting with random linear spaces to reduce down to the isolated case.

A.J. Sommese and C.W. Wampler, Numerical algebraic geometry. *The mathematics of numerical analysis* (Park City, UT 1995). Vol. 32 of *Lectures in Appl. Math.*, 749–763, AMS, Providence, RI, 1996.



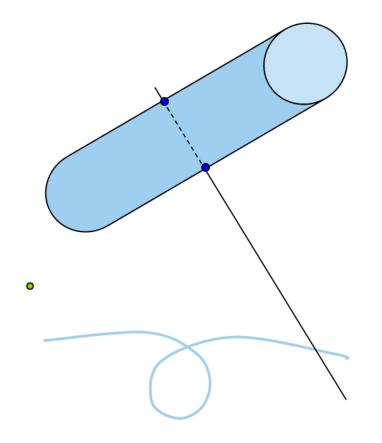






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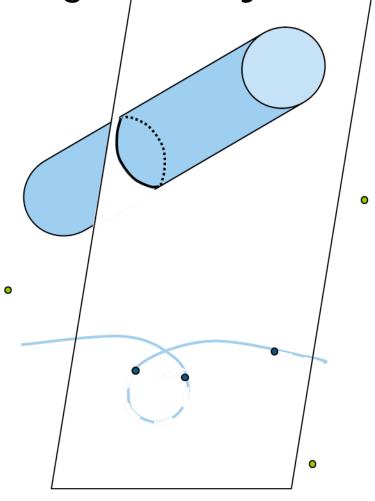






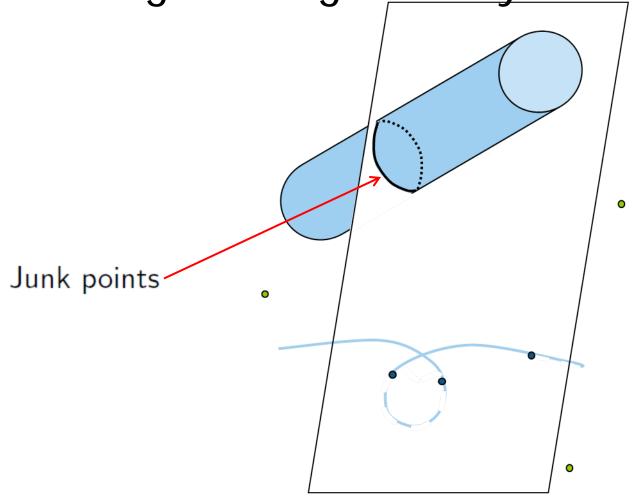
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Compute a local Hilbert function to determine if isolated. D.J. Bates, J.D. Hauenstein, C. Peterson, and A.J. Sommese, A numerical local dimension test for points on the solution set of a system of polynomial equations. SIAM J. Num. Anal., 47(5), 3608–3623, 2009. Junk points





Numerical irreducible decomposition algorithm:

- 1. Compute a witness superset \widehat{W}_k for each dimension k.
 - Compute a superset of the isolated roots of $\begin{vmatrix} f \\ \mathcal{L}_k \end{vmatrix}$.
- 2. Compute a witness set W_k for each k.
 - ▶ Remove the nonisolated roots of $\begin{bmatrix} f \\ \mathcal{L}_k \end{bmatrix}$ from \widehat{W}_k .
- 3. Partition W_k into sets corresponding to the irreducible components of dimension k.





Extension of numerical irreducible decomposition to numerical primary decomposition:

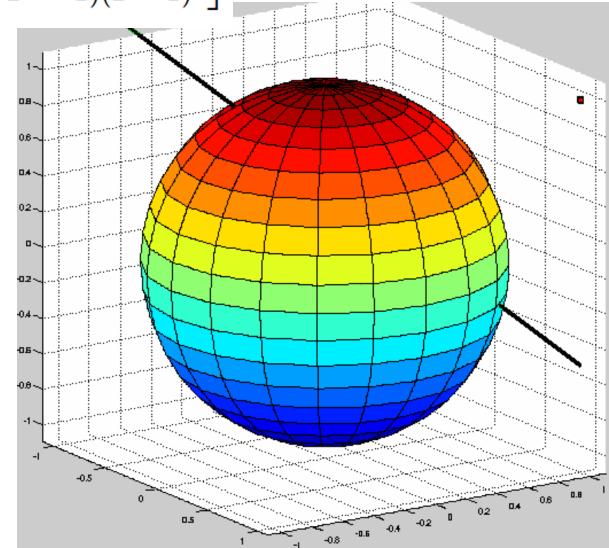
A. Leykin, Numerical primary decomposition. *ISSAC 2008*, 165–172, ACM, New York, 2008.





Example

$$f(x,y,z) = \begin{bmatrix} (x^2 + y^2 + z^2 - 1)(x - 1) \\ (x^2 + y^2 + z^2 - 1)(y - 1) \\ (x^2 + y^2 + z^2 - 1)(z - 1) \end{bmatrix}$$





Example

$$f(x,y,z) = \begin{bmatrix} (x^2 + y^2 + z^2 - 1)(x - 1) \\ (x^2 + y^2 + z^2 - 1)(y - 1) \\ (x^2 + y^2 + z^2 - 1)(z - 1) \end{bmatrix}$$

Dimension 2:

Let L_1, L_2 be random linear polynomials. Solve $\begin{bmatrix} L_1 \\ L_2 \end{bmatrix} = 0$.

$$g_2 = \left[\begin{array}{c} f_1 + \alpha_2 f_2 + \alpha_3 f_3 \\ L_1 \\ L_2 \end{array} \right]$$
 has 2 nonsingular roots that



satisfy f = 0



Example

$$f(x,y,z) = \begin{bmatrix} (x^2 + y^2 + z^2 - 1)(x - 1) \\ (x^2 + y^2 + z^2 - 1)(y - 1) \\ (x^2 + y^2 + z^2 - 1)(z - 1) \end{bmatrix}$$

Dimension 1:

Solve
$$\begin{bmatrix} f \\ L_1 \end{bmatrix} = 0.$$

$$g_1 = \left[egin{array}{c} f_1 + lpha_2 f_2 + lpha_3 f_3 \\ f_2 + eta_3 f_3 \\ L_1 \end{array}
ight]$$
 has 8 singular roots that

satisfy f = 0 - all are junk points.





Example

$$f(x,y,z) = \begin{bmatrix} (x^2 + y^2 + z^2 - 1)(x - 1) \\ (x^2 + y^2 + z^2 - 1)(y - 1) \\ (x^2 + y^2 + z^2 - 1)(z - 1) \end{bmatrix}$$

Dimension 0:

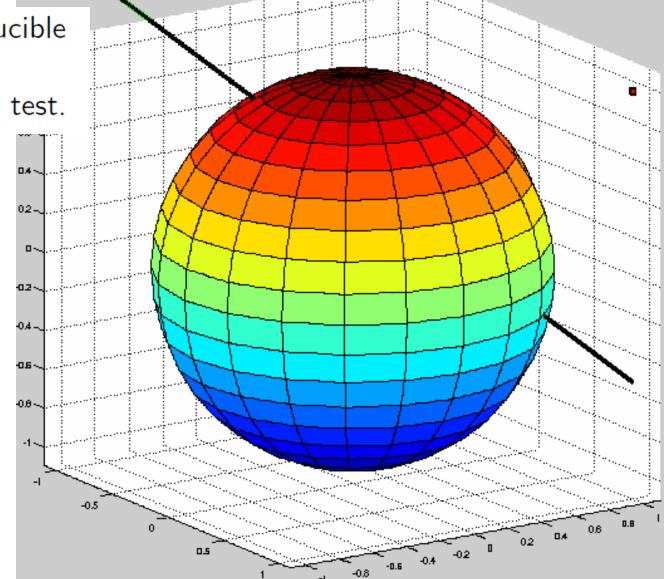
Solve f = 0.

1 nonsingular root and 26 singular roots - all are junk points.

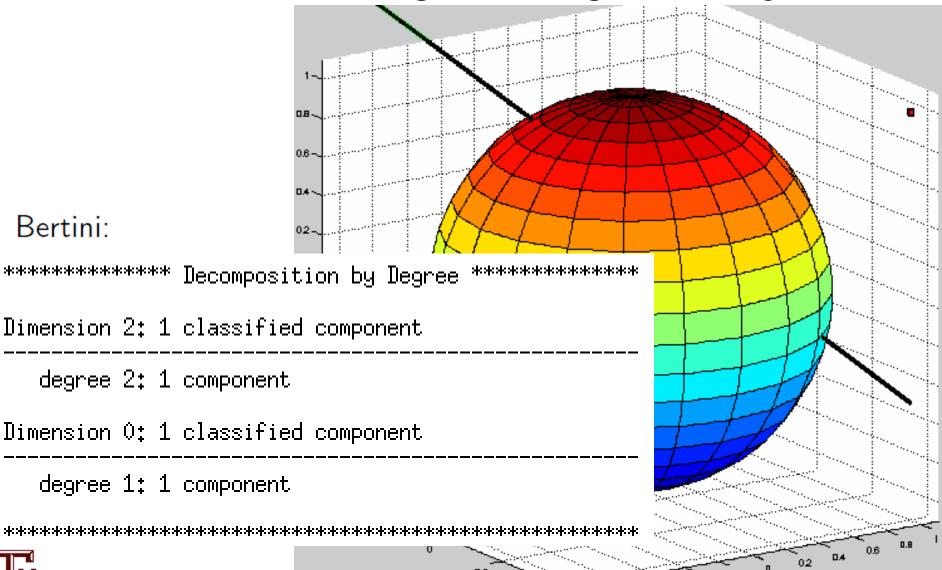




Decompose into irreducible components by using monodromy and trace test.









Bertini:

Given the roots of $\begin{bmatrix} f_1 \\ \vdots \\ f_k \\ L_{k+1} \\ L_{k+2} \\ \vdots \\ L_n \end{bmatrix}$, regeneration $\begin{bmatrix} f_1 \\ \vdots \\ f_k \\ f_{k+1} \\ L_{k+2} \\ \vdots \\ L_n \end{bmatrix}$ computes the roots of $\begin{bmatrix} f_1 \\ \vdots \\ f_k \\ f_{k+1} \\ L_{k+2} \\ \vdots \\ L_n \\ \end{bmatrix}$

$$\begin{bmatrix}
f_1 \\
\vdots \\
f_k \\
f_{k+1} \\
L_{k+2} \\
\vdots \\
L_n
\end{bmatrix}$$

J.D. Hauenstein, A.J. Sommese, and C.W. Wampler, Regeneration homotopies for solving systems of polynomials. To appear in Mathematics of Computation.





Regeneration is effective at computing nonsingular isolated solutions of large scale structured polynomial systems arising in many applications.

The regenerative cascade algorithm applies regeneration to

$$\Re(f) = \begin{bmatrix} f_1 + \alpha_{1,2}f_2 + \alpha_{1,3}f_3 + \dots + \alpha_{1,n}f_n \\ f_2 + \alpha_{2,3}f_3 + \dots + \alpha_{2,n}f_n \\ \vdots \\ f_n \end{bmatrix}$$

in order to compute witness supersets of f.





Step 1:

Move L_{k+1} to $\mathfrak{L}_1, \ldots, \mathfrak{L}_{\deg f_{k+1}}$ using

$$H_i(x,t) = \left[egin{array}{c} f_1 \ dots \ f_k \ dots \ L_{k+2} \ dots \ L_n \end{array}
ight].$$





Step 2:

Introduce f_{k+1} using

$$H(x,t) = \left[egin{array}{c} f_1 \ & dots \ f_k \ f_{k+1}(1-t)+t\prod_{i=1}^{\deg f_{k+1}}\mathfrak{L}_i \ & dots \ L_{k+2} \ & dots \ L_n \end{array}
ight].$$





Adjacent minor system:





Adjacent minor system:

For example:
$$n=3$$

For example:
$$n = 3$$

$$\begin{bmatrix}
 X_1 & X_2 & X_3 \\
 X_4 & X_5 & X_6 \\
 X_7 & X_8 & X_9
\end{bmatrix}$$





Adjacent minor system:

For example:
$$n=3$$

For example:
$$n = 3$$

$$\begin{bmatrix}
 X_1 & X_2 & X_3 \\
 X_4 & X_5 & X_6 \\
 X_7 & X_8 & X_9
\end{bmatrix}$$

$$f_1 = x_1 x_5 - x_2 x_4$$





Adjacent minor system:

For example:
$$n = 3$$

$$f_1 = x_1 x_5 - x_2 x_4$$

$$f_1 = x_1 x_5 - x_2 x_4$$
$$f_2 = x_2 x_6 - x_3 x_5$$





Adjacent minor system:

For example:
$$n = 3$$

$$f_1 = x_1 x_5 - x_2 x_4$$

$$f_3 = x_4 x_8 - x_5 x_7$$

$$f_1 = x_1 x_5 - x_2 x_4$$
$$f_2 = x_2 x_6 - x_3 x_5$$





Adjacent minor system:

Determinants of 2×2 adjacent minors of $3 \times n$ matrix with variable entries.

For example: n = 3

$$f_1 = x_1 x_5 - x_2 x_4$$

$$f_2 = x_2x_6 - x_3x_5$$

$$f_3 = x_4 x_8 - x_5 x_7$$
$$f_4 = x_5 x_9 - x_6 x_8$$

$$f_4 = x_5 x_9 - x_6 x_8$$





Adjacent minor system:

	Membership test			Local dimension test		
n	Regen cascade	Dim-by-dim	Cascade	Regen cascade	Dim-by-dim	Cascade
3	0.1s	0.1s	0.2s	0.1s	0.1s	0.2s
4	0.8s	1.1s	1.3s	0.6s	0.8s	1.1s
5	6.2s	11.9s	11.2s	3.1s	4.6s	7.4s
6	$1 \mathrm{m} 1 \mathrm{s}$	2 m14 s	1 m 34 s	15.6s	29.0s	48.4s
7	10 m 36 s	25 m 39 s	14 m 54 s	$1 \mathrm{m} 16 \mathrm{s}$	3 m 8 s	5m23s
8	2h12m54s	5h21m48s	2h33m5s	6 m 33 s	19m45s	29 m 22 s

J.D. Hauenstein, Regeneration, local dimension, and applications in numerical algebraic geometry. Ph.D. Thesis, University of Notre Dame, Notre Dame, IN, April 2009.





Given
$$A(x) = \begin{bmatrix} a_{1,1}(x) & \cdots & a_{1,n}(x) \\ \vdots & \ddots & \vdots \\ a_{m,1}(x) & \cdots & a_{m,n}(x) \end{bmatrix}$$
, compute the sets

$$S_k(A) = \{x \mid \text{rank } A(x) \leq k\} \text{ and } S_{k,f}(A) = S_k(A) \cap V(f).$$





One way to compute $S_k(A)$ is by creating a polynomial system consisting of the $(k+1) \times (k+1)$ minors of A.

- ► Could yield impractically large system: $\binom{m}{k+1}\binom{n}{k+1}$.
- Each polynomial could consist of many terms.
- Each polynomial could be of high degree.





Let
$$A(x) = \begin{bmatrix} a_{1,1}(x) & \cdots & a_{1,n}(x) \\ \vdots & \ddots & \vdots \\ a_{m,1}(x) & \cdots & a_{m,n}(x) \end{bmatrix}$$
 with $m \ge n$.

Our approach uses the fact that

$$S_k(A) = \{x \mid \text{rank } A(x) \leq k\} = \{x \mid \text{nullity } A(x) \geq n - k\}.$$

D.J. Bates, J.D. Hauenstein, C. Peterson, and A.J. Sommese, Numerical decomposition of the rank-deficiency set of a matrix of multivariate polynomials. *Approximate Commutative Algebra*, edited by L. Robbiano and J. Abbott, *Texts and Monographs in Symbolic Computation*, Springer, 55–77, 2009.





Let
$$\Lambda = \begin{bmatrix} \lambda_{1,1} & \cdots & \lambda_{1,n-k} \\ \vdots & \ddots & \vdots \\ \lambda_{k,1} & \cdots & \lambda_{k,n-k} \end{bmatrix}$$
 and $B \in U(n)$ be random.

We want to solve

$$A(x)B\begin{bmatrix}I_{n-k}\\\Lambda\end{bmatrix}=0.$$

Remarks

- ▶ Added k(n-k) new variables.
- ▶ Consists of m(n k) functions:
 - naturally 2-homogeneous
 - degree in x is same as in A(x)
 - ▶ linear in λ 's
 - straight-line formulation





Example

Compute
$$S_2(A)$$
 for $A = \begin{bmatrix} 0 & a & b & c \\ -a & 0 & d & e \\ -b & -d & 0 & f \\ -c & -e & -f & 0 \end{bmatrix}$.

Determinants: Solve 12 cubics on \mathbb{C}^6 .

Nullity: Solve 8 polynomials of type (1,1) on $\mathbb{C}^6 \times \mathbb{C}^4$.

$$S_2(A) = V(af + cd - be)$$





Example

Compute the singular points of

$$f(x_1, x_2, x_3, x_4) = \begin{bmatrix} x_1 + x_2 + x_3 + x_4 \\ x_1x_2 + x_2x_3 + x_3x_4 + x_4x_1 \\ x_1x_2x_3 + x_2x_3x_4 + x_3x_4x_1 + x_4x_1x_2 \\ x_1x_2x_3x_4 - 1 \end{bmatrix}$$

on the irreducible component $C = \{(x_1, x_2, -x_1, -x_2) \mid x_1x_2 = 1\}.$





Note that rank Jf = 3 generically on C.

Starting with a witness set for C, we compute

$$S_2(Jf)\cap C = \{(1,1,-1,-1),(-1,-1,1,1),(i,-i,-i,i),(-i,i,i,-i)\}.$$



